LOGICS

Exercises on Chapter 2: Sequent Calculus

Symmetries in the system LK

Exercice 1.1 Reversibility

An inference rule $\frac{S_1}{S_2}$ in a formal system S is called *reversible* iff $\frac{S_2}{S_1}$ is a derived rule of the system S. Which rules of system LK are reversible?

Exercice 1.2 Duality

Let us add a connector \top to the language of system LK i.e. the set of connectors is now $\{\bot, \land, \lor, \rightarrow, \neg\}$, the set of quantifiers is $\{\forall, \exists\}$, we have a denumerable set of relation symbols \mathcal{R} and a denumerable set of fonction symbols \mathcal{F} .

To every well-formed formula F, we associate a dual formula D(F) by structural induction:

$$D(\bot) := \top, \quad D(\top) := \bot$$

$$D(R(t_1, ..., t_k)) := R(t_1, ..., t_k) \quad \text{pour tout } R \in \mathcal{R}$$

$$D(F \land G) := D(F) \lor D(G), \quad D(F \lor G) := D(F) \land D(G)$$

$$D(F \to G) := \neg(D(G) \to D(F)), \quad D(\neg F) := \neg D(F)$$

$$D(\forall xF) := \exists x D(F), \quad D(\exists xF) := \forall x D(F)$$

1- Show that, for every formula F, with atomic subformulas A_1, \ldots, A_n , the following equivalence holds:

$$F(\neg A_1, \dots, \neg A_n) \models \neg D(F)(A_1, \dots, A_n).$$

We extend the map $F \mapsto D(F)$ to sequents by:

$$D(F_1,\ldots,F_n \vdash G_1\ldots,G_m) := D(G_1),\ldots,D(G_m) \vdash D(F_1),\ldots,D(F_n).$$

2- Show that, if for every structure \mathcal{T} we have:

$$\mathcal{T} \models S_1 \text{ implies } \mathcal{T} \models S_2$$

then, for every structure \mathcal{T} we also have:

$$\mathcal{T} \models D(S_1) \text{ implies } \mathcal{T} \models D(S_2).$$

- 3- What are the rules of LK, $\frac{S_1}{S_2}$ (where S_1, S_2 are sequents) such that $\frac{D(S_1)}{D(S_2)}$ is still a rule of LK?
- 4- Could you manage to add a rule to LK in such a way that it takes into account the extension to connector \top ? We would like that the sequent $\vdash (\neg \bot \to \top) \land (\top \to \neg \bot)$ be derivable in this extension of LK and we would like to keep (or even improve) the symmetries of system LK.

Examples of derivations

Exercice 1.3 Propositional LK

Find a derivation in LK for the sequents:

Exercice 1.4 LK

Find a derivation in LK for the sequents:

$$\neg \exists x R(x) \models \forall x \neg R(x)$$
$$\neg \forall x R(x) \models \exists x \neg R(x)$$

$$\longmapsto \forall x (Q \lor R(x)) \to Q \lor \forall x R(x)$$

$$-\exists x \forall y (R(y) \rightarrow R(x))$$

Exercice 1.5 Propositional LJ

Find a derivation in LJ for the sequents:

$$\neg (A \lor B) \models \neg A \land \neg B$$

$$\neg A \land \neg B \vdash \neg (A \lor B)$$

$$A \vdash \neg \neg A$$

$$\neg \neg \neg A \vdash \neg A$$

Exercice 1.6 LJ

1- What do you think about the following "derivation"?

$$\begin{array}{lll} R(x) \longmapsto R(x) & \text{ax} \\ \exists x R(x) \longmapsto R(x) & \exists_g \\ \neg R(x), \exists x R(x) \longmapsto & \lnot_g \\ \forall x \neg R(x), \exists x R(x) \longmapsto & \forall_g \\ \forall x \neg R(x) \longmapsto \neg \exists x R(x) & \lnot_d \end{array}$$

2- Does the sequent $\forall x \neg R(x) \models \neg \exists x R(x)$ admit a derivation in LJ?

Properties of derivations

Exercice 1.7 Right-elimination

Let us consider the following right-elimination rule:

$$\frac{\Gamma \longmapsto A \wedge B}{\Gamma \longmapsto A}$$

(where Γ is a multi-set of formulas and A, B are formulas).

- 1- Is this rule a derived rule of LK?
- 2- Is this rule a derived rule of LK \ {cut} ?

Exercice 1.8 Contractions

Let us consider the sequent $S := \exists x (R(a) \lor R(b) \to R(x)).$

- 1- Find a derivation of S in LK.
- 2- Find a cut-free derivation of S in LK.
- 3- Prove that there does not exist any derivation, without neither cut nor contraction, of S in LK
- 4- Does there exist some derivation, without contraction, of S in LK?